Note Title 9/5/2007

Digital Systems and Information Representation

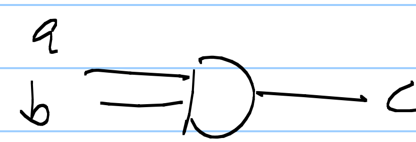
Base signals two-valued 0 or 1

"bit" smallest unambiguous unit of information

propositional calc.

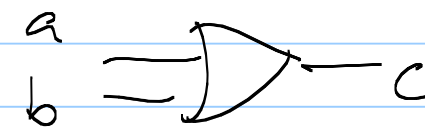
operations

AND $a \cdot b = c$



c is true iff a is true and b is true

OR $a+b=c$

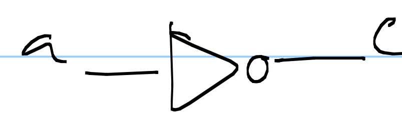


c iff a or b or both

NOT $a'=c$

$\bar{a}=c$

a is true iff c is false



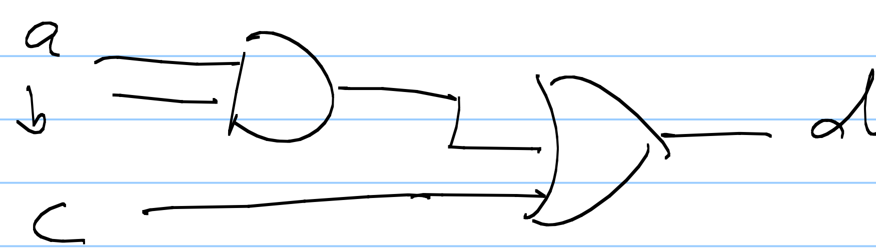
$\sim a = c$

$\neg a = c$

e.g.

a	passed	microeconomics
b	"	macro "
c	"	econ survey
d	met	econ grad. req.

$d = a \cdot b + c$



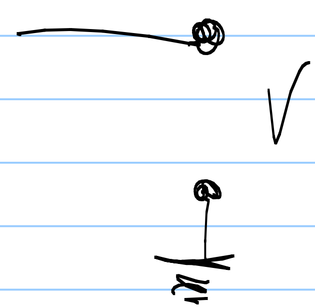
1 true

0 false

Boolean algebra

physical representation

Wire



$V > V_H \Rightarrow 1$

$V < V_L \Rightarrow 0$

$V_H > V_L$

region between V_L and V_H is undefined

two values are not enough

→ numbers, text, images, audio, video, ...

sequence of number systems

1. "counting" numbers 1, 2, 3, . . (positive integers)
2. What if I don't have any stuff?
   - add 0 as a number

· what if I take away more than I have?

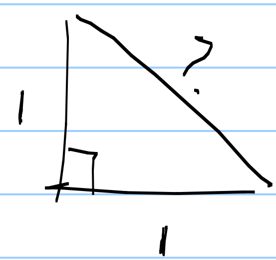
new concept: negative numbers

can solve equation $x + a = 0$ integer a

Why can't I divide my stuff 3 ways?

incorporate rational numbers

eqn $ax + b = 0$ integers a, b if $a \neq 0$



how long is this line?

expand to irrational numbers

and get "reals"

$$x^2 - 2 = 0$$

what about $x^2 + 2 = 0$

complex numbers

vector number system with 2 components

(a, b) a, b both "real" numbers

rules

equality $(a, b) = (c, d)$

iff $a = c$ and $b = d$

add

$(a,b) + (c,d) = (a+c, b+d)$

mult.

2 complex #s

$(a,b) \times (c,d) = (ac - bd, ad + bc)$

scalar and complex

$c \times (a,b) = (ac, bc)$

complex #s w/ 2^{nd} comp 0 have same properties as reals

$$(a, 0) = (c, 0) \text{ iff } a = c$$

$$(a, 0) + (c, 0) = (a+c, 0)$$

$$(a, 0) \times (c, 0) = (ac, 0)$$

$$(a, 0) \times c = (ac, 0)$$

1st comp. of complex is called "real"

then follow historical naming conventions rational → irrational

real → imaginary

2^{nd} component of complex # known as imaginary component

equ. $x^2 + 2 = 0$

$$x^2 + (2, 0) = (0, 0)$$

insert $x = (0, \sqrt{2})$

$x^2 + (2,0) = (0,0)$

$(0,\sqrt{2}) \cdot (0,\sqrt{2}) + (2,0) = (0,0)$

$(-2,0) + (2,0) = (0,0)$ ✓

$x^2 + 1 = 0$

$x^2 = -1$

$x = \sqrt{-1}$

$x = (0,1)$

$(0,1)^2 + (1,0) = (0,0)$

$(-1,0) + (1,0) = (0,0)$ ✓

$x = \sqrt{-1} = (0,1)$

new symbol $i \triangleq \sqrt{-1} = (0,1)$

(a,b) can now be written $a + ib$

$$(a,b) = a + ib \qquad i = (0,1)$$

$$a + ib = (a,0) + (0,1) \times (b,0)$$
$$= (a,0) + (0,b)$$
$$= (a,b)$$

EEs use j instead of i

can now solve

$$a_0 x^n + a_1 x^{n-1} + a_2 x^{n-2} \cdots a_{n-1} x + a_n = 0$$

if a's complex, $n \geq 1$ and $a_0 \neq 0 \Rightarrow n$ roots

Back to representing signal (numbers) that are multi-valued

positional number systems

base 10 $\quad xyz_{10} = x \cdot 10^2 + y \cdot 10^1 + z \cdot 10^0$

x, y, z have values 0 to 9

base 2 $\quad xyz_2 = x \cdot 2^2 + y \cdot 2^1 + z \cdot 2^0$

x, y, z value values 0 or 1

x	y	z	
0	0	0	0
0	0	1	1
0	1	0	2
0	1	1	3
1	0	0	4
1	0	1	5
1	1	0	6
1	1	1	7

base 2 called binary

with fixed # of bits one can rep negative numbers multiple ways fixed # of digits (bits)

e.g. 4-bit binary # system

unsigned range 0 to 15

0000 to 1111

w/ n bits ⇒ range 0 to $2^n - 1$

offset or bias (e.g. -7)

range -7 to +8

sign-magnitude

1st bit sign (0 = positive, 1 = negative)

bits 2, 3, 4 magnitude range 0 to 7

overall range -7 to 7

what about 1000? -0

w/ n bits $-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

2's complement (radix complement)

$$wxyz = w \cdot -(2)^3 + x \cdot 2^2 + y \cdot 2^1 + z \cdot 2^0$$

$$= w \cdot (-8) + x \cdot 4 \cdot y \cdot 2 + z$$

range is -8 to $+7$

1st bit still sign bit

only one zero 0000

positional # system

negative weight on msb

n bits

range $-(2^{n-1})$ to $+(2^{n-1}-1)$

only need adder for add/sub

2's complement is almost universal